

# Automatic Differentiation – Lecture No 2

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Computing with the C/C++ `single` format

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Example 1: Repeated addition

$$\sum_{i=1}^{10^3} \langle 10^{-3} \rangle = 0.999990701675415,$$

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Example 2: Order of summation

$$1 + \frac{1}{2} + \frac{1}{3} + \cdots + \frac{1}{10^6} = 14.357357,$$

$$\frac{1}{10^6} + \cdots + \frac{1}{3} + \frac{1}{2} + 1 = 14.392651.$$

# Are floating point computations reliable?

Given the point  $(x, y) = (77617, 33096)$ , evaluate the function

$$f(x, y) = 333.75y^6 + x^2(11x^2y^2 - y^6 - 121y^4 - 2) + 5.5y^8 + x/(2y)$$

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IBM S/370 ( $\beta = 16$ ) with FORTRAN:

type	$p$	$f(x, y)$
REAL*4	24	1.172603 ...
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Pentium III ( $\beta = 2$ ) with C/C++ (gcc/g++):

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Correct answer:  $-0.8273960599 \dots$



# Are integer computations reliable?

## Example 3: The harmonic series

If we define

$$S_N = \sum_{k=1}^N \frac{1}{k},$$

then  $\lim_{N \rightarrow \infty} S_N = +\infty$ . The computer also gets this result, but for the entirely wrong reason (integer wrapping).

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## Example 4: Elementary Taylor series

Now define

$$S_N = \sum_{k=0}^N \frac{1}{k!},$$

then  $\lim_{N \rightarrow \infty} S_N = e \approx 2.7182818$ . The integer wrapping produces a sequence that is *not* strictly increasing:

# Are integer computations reliable?

S_0 = 1.0000000000000000	S_18 = 2.718281826004540	S
S_1 = 2.0000000000000000	S_19 = 2.718281835125155	
S_2 = 2.5000000000000000	S_20 = 2.718281834649448	S
S_3 = 2.6666666666666667	S_21 = 2.718281833812708	S
S_4 = 2.7083333333333333	S_22 = 2.718281831899620	S
S_5 = 2.7166666666666666	S_23 = 2.718281833059102	
S_6 = 2.7180555555555555	S_24 = 2.718281831770353	S
S_7 = 2.718253968253968	S_25 = 2.718281832252007	
S_8 = 2.718278769841270	S_26 = 2.718281831712599	S
S_9 = 2.718281525573192	S_27 = 2.718281832386097	
S_10 = 2.718281801146385	S_28 = 2.718281831659211	S
S_11 = 2.718281826198493	S_29 = 2.718281830853743	S
S_12 = 2.718281828286169	S_30 = 2.718281831563322	
S_13 = 2.718281828803753	S_31 = 2.718281832917973	
S_14 = 2.718281829585647	S_32 = 2.718281832452312	S
S_15 = 2.718281830084572	S_33 = 2.718281831986650	S
S_16 = 2.718281830583527	S_34 =	inf
S_17 = 2.718281827117590		S

# How do we control rounding errors?

Round each partial result both ways

If  $x, y \in \mathbb{F}$  and  $\star \in \{+, -, \times, \div\}$ , we can enclose the exact result in an *interval*:

$$x \star y \in [\nabla(x \star y), \Delta(x \star y)].$$



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Question

How do we compute with intervals? And why, really?

## Definition

If  $\star$  is one of the operators  $+$ ,  $-$ ,  $\times$ ,  $\div$ , and if  $\mathbb{A}, \mathbb{B} \in \mathbb{R}$ , then

$$\mathbb{A} \star \mathbb{B} = \{a \star b : a \in \mathbb{A}, b \in \mathbb{B}\},$$

except that  $\mathbb{A} \div \mathbb{B}$  is undefined if  $0 \in \mathbb{B}$ .



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## Simple arithmetic

$$\mathbb{A} + \mathbb{B} = [\underline{\mathbb{A}} + \underline{\mathbb{B}}, \overline{\mathbb{A}} + \overline{\mathbb{B}}]$$

$$\mathbb{A} - \mathbb{B} = [\underline{\mathbb{A}} - \overline{\mathbb{B}}, \overline{\mathbb{A}} - \underline{\mathbb{B}}]$$

$$\mathbb{A} \times \mathbb{B} = [\min\{\underline{\mathbb{A}}\underline{\mathbb{B}}, \underline{\mathbb{A}}\overline{\mathbb{B}}, \overline{\mathbb{A}}\underline{\mathbb{B}}, \overline{\mathbb{A}}\overline{\mathbb{B}}\}, \max\{\underline{\mathbb{A}}\underline{\mathbb{B}}, \underline{\mathbb{A}}\overline{\mathbb{B}}, \overline{\mathbb{A}}\underline{\mathbb{B}}, \overline{\mathbb{A}}\overline{\mathbb{B}}\}]$$

$$\mathbb{A} \div \mathbb{B} = \mathbb{A} \times [1/\overline{\mathbb{B}}, 1/\underline{\mathbb{B}}], \quad \text{if } 0 \notin \mathbb{B}.$$





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On a computer we use *directed rounding*, e.g.

$$\mathbb{A} + \mathbb{B} = [\nabla(\underline{\mathbb{A}} \oplus \underline{\mathbb{B}}), \Delta(\overline{\mathbb{A}} \oplus \overline{\mathbb{B}})].$$

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# Interval arithmetic - implementations

It is natural to implement the `interval` class as of two numbers – the endpoints of the interval:

```
01 function iv = interval(lo, hi)
02 % A naive interval class constructor.
03 if nargin == 1
04     hi = lo;
05 elseif ( hi < lo )
06     error('The endpoints do not define an interval.');
```

07 end

```
08 iv.lo = lo; iv.hi = hi;
09 iv = class(iv, 'interval');
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08 iv.lo = lo; iv.hi = hi;
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By including lines 03 and 04, we allow the constructor to automatically cast a single number  $x$  to a thin interval.

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```
01 function display(iv)
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```

We can now input/output intervals within the MATLAB environment:

```
>> a = interval(3, 4), b = interval(2, 5), c = interval(1)
a =
 [3.000000000000000000, 4.000000000000000000]
b =
 [2.000000000000000000, 5.000000000000000000]
c =
 [1.000000000000000000, 1.000000000000000000]
```



# Interval arithmetic - implementations

Implementing the arithmetic operations is straight-forward:



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```
01 function result = plus(a, b)
02 % Overloading the '+' operator for intervals.
03 [a, b] = cast(a, b);
04 setround(-inf);
05 lo = a.lo + b.lo;
06 setround(+inf);
07 hi = a.hi + b.hi;
08 setround(0.5);
09 result = interval(lo, hi);
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Note the call to `setround` on lines 04, 06, and 08:

```
01 function setround(rnd)
02 % A switch for changing rounding mode. The arguments
03 % {+inf, -inf, 0.5, 0} correspond to the roundings
04 % {upward, downward, to nearest, to zero}, respectively.
05 system_dependent('setround', rnd);
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This is a hidden (undocumented) feature of MATLAB.

# Interval arithmetic - implementations

Performing some simple interval calculations, we have:

```
>> a+b, a-b, a*b, a/b
```

```
ans =
```

```
[5.000000000000000000, 9.000000000000000000]
```

```
ans =
```

```
[-2.000000000000000000, 2.000000000000000000]
```

```
ans =
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[6.000000000000000000, 20.000000000000000000]
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```
ans =
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[0.599999999999999998, 2.000000000000000000]
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```

Note the outward rounding in the left endpoint of the last result. Now let's try to find the smallest interval containing  $1/10$ .

```
>> interval(1/10)
ans =
  [0.10000000000000001, 0.10000000000000001]
>> interval(1)/10
ans =
  [0.09999999999999999, 0.10000000000000001]
```

## Range enclosure

Extend a real-valued function  $f$  to an interval-valued  $F$ :

$$R(f; \mathbb{X}) = \{f(x) : x \in \mathbb{X}\} \subseteq F(\mathbb{X})$$

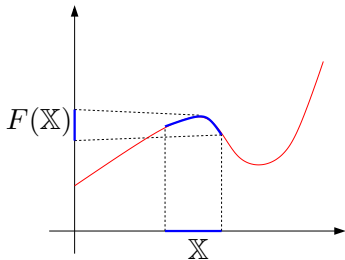
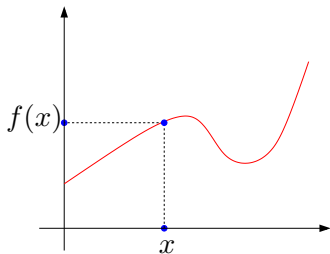


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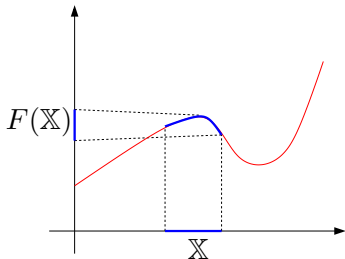
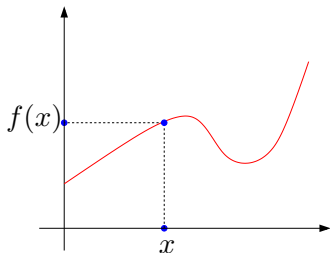


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$y \notin F(\mathbb{X})$  implies that  $f(x) \neq y$  for all  $x \in \mathbb{X}$ .

Some explicit formulas are given below:

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$$\begin{aligned} e^{\underline{X}} &= [e^{\underline{X}}, e^{\overline{X}}] \\ \sqrt{\underline{X}} &= [\sqrt{\underline{X}}, \sqrt{\overline{X}}] && \text{if } 0 \leq \underline{X} \\ \log \underline{X} &= [\log \underline{X}, \log \overline{X}] && \text{if } 0 < \underline{X} \\ \arctan \underline{X} &= [\arctan \underline{X}, \arctan \overline{X}] \quad . \end{aligned}$$



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Set  $S^+ = \{2k\pi + \pi/2: k \in \mathbb{Z}\}$  and  $S^- = \{2k\pi - \pi/2: k \in \mathbb{Z}\}$ .  
Then  $\sin \mathbb{X}$  is given by

$$\begin{cases} [-1, 1] & : \text{if } \mathbb{X} \cap S^- \neq \emptyset \text{ and } \mathbb{X} \cap S^+ \neq \emptyset, \\ [-1, \max\{\sin \underline{\mathbb{X}}, \sin \overline{\mathbb{X}}\}] & : \text{if } \mathbb{X} \cap S^- \neq \emptyset \text{ and } \mathbb{X} \cap S^+ = \emptyset, \\ [\min\{\sin \underline{\mathbb{X}}, \sin \overline{\mathbb{X}}\}, 1] & : \text{if } \mathbb{X} \cap S^- = \emptyset \text{ and } \mathbb{X} \cap S^+ \neq \emptyset, \\ [\min\{\sin \underline{\mathbb{X}}, \sin \overline{\mathbb{X}}\}, \max\{\sin \underline{\mathbb{X}}, \sin \overline{\mathbb{X}}\}] & : \text{if } \mathbb{X} \cap S^- = \emptyset \text{ and } \mathbb{X} \cap S^+ = \emptyset. \end{cases}$$

# Interval-valued functions

A simple (and incorrect) implementation of sin is the following:

```
01 function result = sin(x)
02 % Overloading the 'sin' operator.
03 Sp = (2*floor(x.hi/(2*pi) - 1/4)*pi + pi/2 <= x);
04 Sm = (2*floor(x.hi/(2*pi) + 1/4)*pi - pi/2 <= x);
05 system_dependent('setround',-inf);
06 min_sin = min(sin(x.lo),sin(x.hi));
07 system_dependent('setround',+inf);
08 max_sin = max(sin(x.lo),sin(x.hi));
09 system_dependent('setround',0.5);
10 if ( Sm && Sp )
11     result = interval(-1,+1);
12 elseif Sm
13     result = interval(-1, max_sin);
14 elseif Sp
15     result = interval(min_sin, +1);
16 else
17     result = interval(min_sin, max_sin);
18 end
```



## Exercise

Draw an accurate graph of the function  $f(x) = \cos^3 x + \sin x$  over the domain  $[-5, 5]$ .



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## Solution

Define  $F(\mathbb{X}) = \cos^3 \mathbb{X} + \sin \mathbb{X}$ , and adaptively bisect the domain  $\mathbb{X}_0 = [-5, 5]$  into smaller pieces  $\mathbb{X}_0 = \cup_{i=1}^N \mathbb{X}_i$  until we arrive at some desired accuracy, e.g.  $\max_i \text{width}(F(\mathbb{X}_i)) \leq \text{TOL}$ .



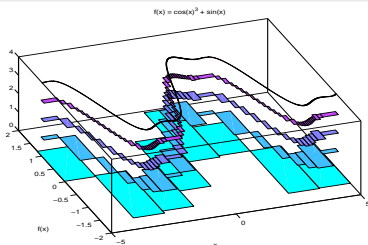
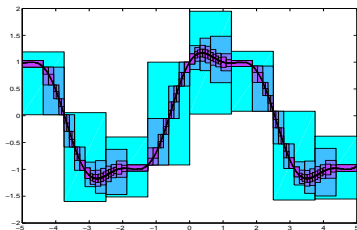
# Graph Enclosures

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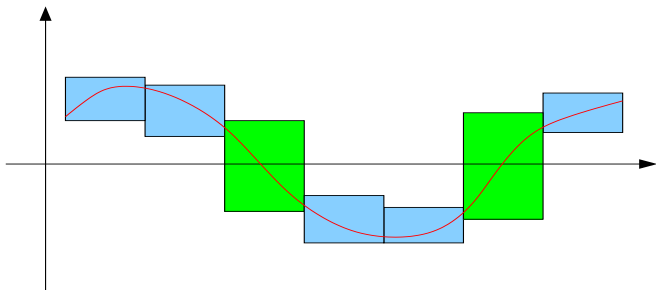
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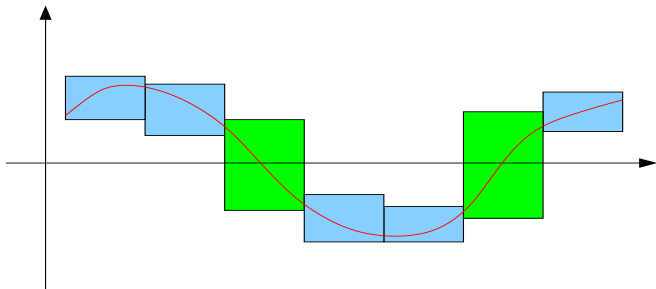




# Solving non-linear equations

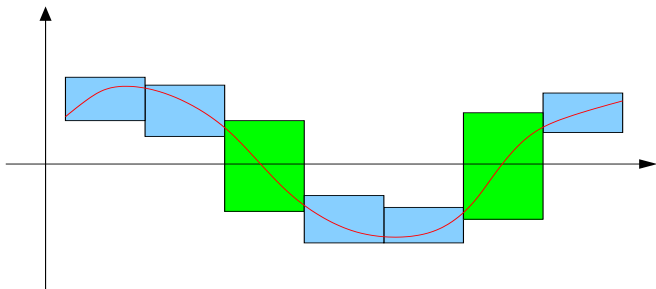


# Solving non-linear equations



*Consider everything. Keep what is good.  
Avoid evil whenever you recognize it.*  
St. Paul, ca. 50 A.D. (The Bible, 1 Thess. 5:21-22)

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No solutions can be missed!

# Solving non-linear equation

The code is transparent and natural

```
01 function bisect(fcnName, X, tol)
02 f = inline(fcnName);
03 if ( 0 <= f(X) )           % If f(X) contains zero...
04     if Diam(X) < tol      % and the tolerance is met...
05         X                 % print the interval X.
06     else                  % Otherwise, divide and conquer.
07         bisect(fcnName, interval(Inf(X), Mid(X)), tol);
08         bisect(fcnName, interval(Mid(X), Sup(X)), tol);
09     end
10 end
```



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```

## Nice property

If  $F$  is well-defined on the domain, the algorithm produces an enclosure of *all* zeros of  $f$ . [No existence is established, however.]



# Existence and uniqueness

Existence and uniqueness require *fixed point* theorems.

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## Brouwer's fixed point theorem

Let  $B$  be homeomorphic to the closed unit ball in  $\mathbb{R}^n$ . Then given any continuous mapping  $f: B \rightarrow B$  there exists  $x \in B$  such that  $f(x) = x$ .



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## Schauder's fixed point theorem

Let  $X$  be a normed vector space, and let  $K \subset X$  be a non-empty, compact, and convex set. Then given any continuous mapping  $f: K \rightarrow K$  there exists  $x \in K$  such that  $f(x) = x$ .



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## Brouwer's fixed point theorem

Let  $B$  be homeomorphic to the closed unit ball in  $\mathbb{R}^n$ . Then given any continuous mapping  $f: B \rightarrow B$  there exists  $x \in B$  such that  $f(x) = x$ .

## Schauder's fixed point theorem

Let  $X$  be a normed vector space, and let  $K \subset X$  be a non-empty, compact, and convex set. Then given any continuous mapping  $f: K \rightarrow K$  there exists  $x \in K$  such that  $f(x) = x$ .

## Banach's fixed point theorem

If  $f$  is a contraction defined on a complete metric space  $X$ , then there exists a unique  $x \in X$  such that  $f(x) = x$ .

## Theorem

Let  $f \in C^1(\mathbb{R}, \mathbb{R})$ , and set  $\check{x} = \text{mid}(\mathbb{X})$ . We define

$$N_f(\mathbb{X}) \stackrel{\text{def}}{=} N_f(\mathbb{X}, \check{x}) = \check{x} - f(\check{x})/F'(\mathbb{X}).$$

If  $N_f(\mathbb{X})$  is well-defined, then the following statements hold:

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## Proof.

- (1) Follows from the MVT;
- (2) The contra-positive statement of (1);
- (3) Existence from Brouwer's fixed point theorem;  
Uniqueness from non-vanishing  $f'$ .

## Algorithm

Starting from an initial search region  $\mathbb{X}_0$ , we form the sequence

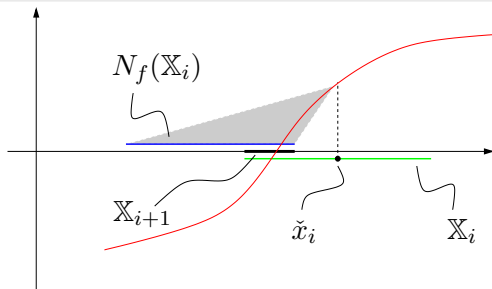
$$\mathbb{X}_{i+1} = N_f(\mathbb{X}_i) \cap \mathbb{X}_i \quad i = 0, 1, \dots$$



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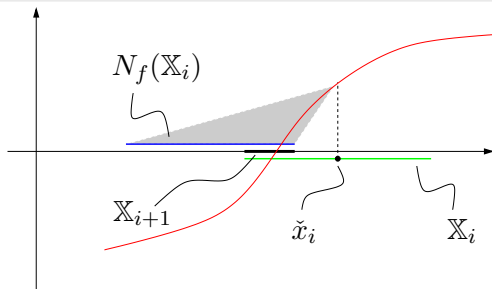




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## Performance

If well-defined, this method is never worse than bisection, and it converges quadratically fast under mild conditions.

## Example

Let  $f(x) = -2.001 + 3x - x^3$ , and  $\mathbb{X}_0 = [-3, -3/2]$ . Then  $F'(\mathbb{X}_0) = [-24, -15/4]$ , so  $N_f(\mathbb{X}_0)$  is well-defined, and the above theorem holds.



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X(0) = [-3.000000000000000, -1.500000000000000]; rad = 7.50000e-01

X(1) = [-2.140015625000001, -1.546099999999999]; rad = 2.96958e-01

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Unique root in  $-2.00011110288172 \pm 1.555e-15$

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## Question:

What do we do when  $\mathbb{X}_0$  contains several zeros?

# The Krawczyk method

If  $f$  has a zero  $x^*$  in  $\mathbb{X}$ , then for any  $x \in \mathbb{X}$ , we can enclose the zero via

$$x^* \in x - Cf(x) - (1 - CF'(\mathbb{X}))(x - \mathbb{X}) \stackrel{\text{def}}{=} K_f(\mathbb{X}, x, C).$$



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Good choices are  $x = \check{x}$  and  $C = 1/f'(\check{x})$ , yielding the *Krawczyk operator*

$$K_f(\mathbb{X}) \stackrel{\text{def}}{=} \check{x} - \frac{f(\check{x})}{f'(\check{x})} - \left(1 - \frac{F'(\mathbb{X})}{f'(\check{x})}\right) [-r, r],$$

where we use the notation  $r = \text{rad}(\mathbb{X})$ .

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## Theorem

Assume that  $K_f(\mathbb{X})$  is well-defined. Then the following statements hold:

- (1) if  $\mathbb{X}$  contains a zero  $x^*$  of  $f$ , then so does  $K_f(\mathbb{X}) \cap \mathbb{X}$ ;
- (2) if  $K_f(\mathbb{X}) \cap \mathbb{X} = \emptyset$ , then  $\mathbb{X}$  contains no zeros of  $f$ ;
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## Example

Let  $f(x) = \sin x(x - \cos x)$ , and  $\mathbb{X}_0 = [1, 15]$ .



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Tolerance : 1e-13

Function calls : 59

Unique zero in the interval 3.14159265358979 [08,76]

Unique zero in the interval 6.28318530717958 [44,89]

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Of course, all derivatives are computed using AD-techniques, overloaded with intervals.

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Compute (approximate/enclose) the definite integral

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## Naive approach:

Split the domain of integration into  $N$  equally wide subintervals: we set  $h = (b - a)/N$  and  $x_i = a + ih$ ,  $i = 0, \dots, N$ , and enclose the integrand via IA.



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This produces the enclosure

$$\int_a^b f(x)dx \in I_a^b(f, N) \stackrel{\text{def}}{=} h \sum_{i=1}^N F([x_{i-1}, x_i]),$$

which satisfies  $w(I_a^b(f, N)) = \mathcal{O}(1/N)$ .

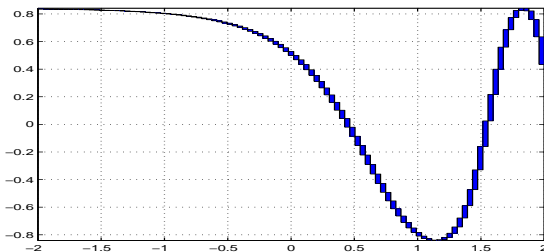
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Enclose the definite integral  $\int_{-2}^2 \sin(\cos(e^x))dx$ .

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$N$	$I_{-2}^2(f, N)$	$w(I_{-2}^2(f, N))$
$10^0$	$[-3.36588, 3.36588]$	$6.73177 \cdot 10^0$
$10^2$	$[1.26250, 1.41323]$	$1.50729 \cdot 10^{-1}$
$10^4$	$[1.33791, 1.33942]$	$1.50756 \cdot 10^{-3}$
$10^6$	$[1.33866, 1.33868]$	$1.50758 \cdot 10^{-5}$





Taylor series approach:

Generate tighter bounds on the integrand via AD.

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For all  $x \in \mathbb{X}$ , we have

$$\begin{aligned} f(x) &= \sum_{k=0}^{n-1} f_k(\tilde{x})(x - \tilde{x})^k + f_n(\zeta_x)(x - \tilde{x})^n \\ &\in \sum_{k=0}^n f_k(\tilde{x})(x - \tilde{x})^k + [-\varepsilon_n, \varepsilon_n]|x - \tilde{x}|^n, \end{aligned}$$

where  $\varepsilon_n = \text{mag}(F_n(\mathbb{X}) - f_n(\tilde{x}))$ .



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where  $\varepsilon_n = \text{mag}(F_n(\mathbb{X}) - f_n(\tilde{x}))$ .

We are now prepared to compute the integral itself:

$$\begin{aligned} \int_{\tilde{x}-r}^{\tilde{x}+r} f(x) dx &\in \int_{\tilde{x}-r}^{\tilde{x}+r} \left( \sum_{k=0}^n f_k(\tilde{x})(x - \tilde{x})^k + [-\varepsilon_n, \varepsilon_n]|x - \tilde{x}|^n \right) dx \\ &= \sum_{k=0}^n f_k(\tilde{x}) \int_{-r}^r x^k dx + [-\varepsilon_n, \varepsilon_n] \int_{-r}^r |x|^n dx. \end{aligned}$$

Continuing the calculation, we see a lot of cancellation:

$$\begin{aligned}\int_{\check{x}-r}^{\check{x}+r} f(x) dx &\in \sum_{k=0}^n f_k(\check{x}) \int_{-r}^r x^k dx + [-\varepsilon_n, \varepsilon_n] \int_{-r}^r |x|^n dx \\ &= \sum_{k=0}^{\lfloor n/2 \rfloor} f_{2k}(\check{x}) \int_{-r}^r x^{2k} dx + [-\varepsilon_n, \varepsilon_n] \int_{-r}^r |x|^n dx \\ &= 2 \left( \sum_{k=0}^{\lfloor n/2 \rfloor} f_{2k}(\check{x}) \frac{r^{2k+1}}{2k+1} + [-\varepsilon_n, \varepsilon_n] \frac{r^{n+1}}{n+1} \right).\end{aligned}$$



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 \end{aligned}$$

Partition the domain of integration:  $a = x_0 < x_1 < \dots < x_N = b$ .

$$\begin{aligned}
 \int_a^b f(x) dx &= \sum_{i=1}^N \int_{x_{i-1}}^{x_i} f(x) dx = \sum_{i=1}^N \int_{\check{x}_i-r_i}^{\check{x}_i+r_i} f(x) dx \\
 &\in 2 \sum_{i=1}^N \left( \sum_{k=0}^{\lfloor n/2 \rfloor} f_{2k}(\check{x}_i) \frac{r_i^{2k+1}}{2k+1} + [-\varepsilon_{n,i}, \varepsilon_{n,i}] \frac{r_i^{n+1}}{n+1} \right).
 \end{aligned}$$

Uniform partition:

$N$	$E_{-2}^2(f, 6, N)$	$w(E_{-2}^2(f, 6, N))$
9	[0.86325178469, 1.81128961988]	$9.4804 \cdot 10^{-1}$
12	[1.28416304745, 1.39316025451]	$1.0900 \cdot 10^{-1}$
21	[1.33783795371, 1.33950680633]	$1.6689 \cdot 10^{-3}$
75	[1.33866863493, 1.33866878008]	$1.4514 \cdot 10^{-7}$



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## Adaptive partition:

TOL	$A_{-2}^2(f, 6, \text{TOL})$	$w(A_{-2}^2(f, 6, \text{TOL}))$	$N_{\text{TOL}}$
$10^{-1}$	[1.33229594606, 1.34500942603]	$1.2713 \cdot 10^{-2}$	9
$10^{-2}$	[1.33822575109, 1.33911045235]	$8.8470 \cdot 10^{-4}$	12
$10^{-4}$	[1.33866170207, 1.33867571626]	$1.4014 \cdot 10^{-5}$	21
$10^{-8}$	[1.33866870618, 1.33866870862]	$2.4304 \cdot 10^{-9}$	75



## Example (A bonus problem)

Compute the integral  $\int_0^8 \sin(x + e^x) dx$ .



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A regular MATLAB session:

```
% Define the integrand and domain.  
>> f = vectorize(inline('sin(x + exp(x))'));  
>> a = 0; b = 8;  
% Compute the integral using MATLAB's 'quad'.  
>> q = quad(f,a,b)  
q =  
    0.251102722027180
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Using the adaptive validated integrator:

```
$$ ./adQuad 0 8 20 1e-10  
Partitions: 874  
CPU time   : 0.45 seconds  
Integral   : 0.3474001726 [492276, 652638]
```

## Interval Computations Web Page

<http://www.cs.utep.edu/interval-comp>



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## **PROFIL/BIAS**

<http://www.ti3.tu-harburg.de/Software/PROFIEnglisch.html>

